

# Mean field for large bike sharing systems

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# Motivation

- ▶ Bike sharing systems become popular
  - ▶ Since **Velib** in Paris 2007
  - ▶ More than 200 programs around the world  
Ex: Barcelona, Lyon, London, Montreal, Washington, Hangzou (China)
- ▶ why?
  - ▶ Good for the town (pollution, traffic jams, health);
  - ▶ Good for the citizen (not to buy, to park the bike, no risk of theft).
- ▶ **BUT**

Empty station



:(

Full station



:(

Good stations



:)

**problematic stations**

Empty station



:(

Full station



:(

Good stations



:)

**problematic stations**

# Le Monde

Mercredi 13 Juillet 2011

& vous

**MODE DE VIE**

## **Le Vélib' ne se trouve pas toujours à sa place**

**Quatre ans après son lancement, le système se heurte encore au problème des bornes vides ou engorgées**

# Velib : a large stochastic network



Map of Velib' stations in Paris.

- ▶ **1500** stations
- ▶ **20000** bikes

## Usage:

- ▶ Take a bike from any station.
- ▶ Use it.
- ▶ Return it to a station of your choice

- ▶ **Performance:** low proportion of
  - ▶ empty stations
  - ▶ full stations
- ▶ huge number of nodes (stations)  $\rightarrow +\infty$   
 $\Rightarrow$  mean field

## A simple case: homogeneous

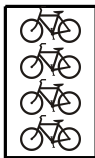
**$C = 4$**



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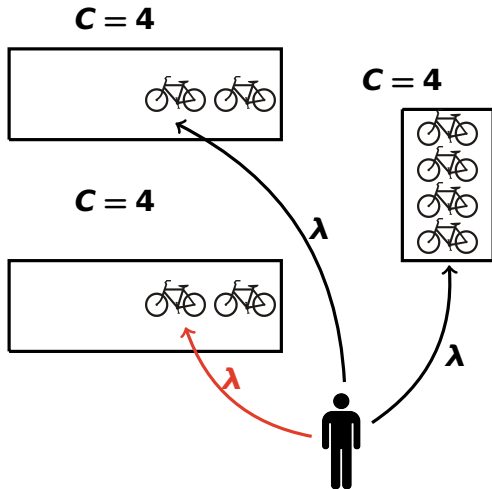
For all  **$N$**  stations:

- Capacity  **$C$**

Extended to non-homogeneous case

- arrival rate, routing probability

## A simple case: homogeneous



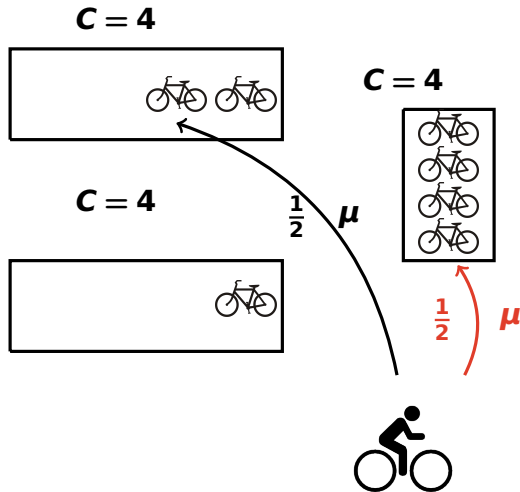
For all  $N$  stations:

- Capacity  $C$
- Arrival rate  $\lambda$

Extended to non-homogeneous case

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## A simple case: homogeneous



For all  $N$  stations:

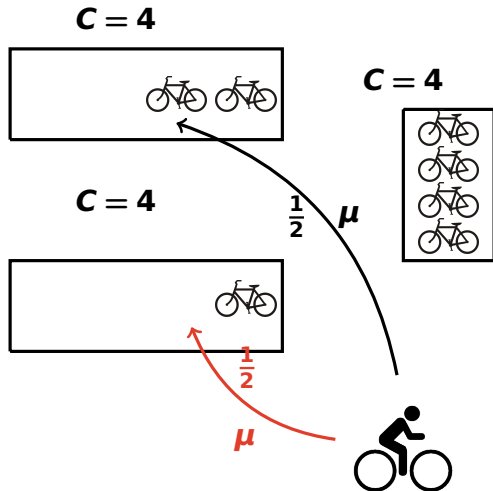
- Capacity  $C$
- Arrival rate  $\lambda$
- Travel time: exponential of mean  $1/\mu$
- Choose one station at random.

Extended to non-homogeneous case

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## A simple case: homogeneous



For all  $N$  stations:

- ▶ Capacity  $C$
- ▶ Arrival rate  $\lambda$
- ▶ Travel time: exponential of mean  $1/\mu$
- ▶ Choose one station at random.
- ▶ If full, try again at random ( $\approx$  local search)

Extended to non-homogeneous case

- ▶ arrival rate, routing probability

## A first result

The stationary distribution of the **number of bikes in a station** converges to ***geom*( $\rho$ )** on  **$0, \dots, C$** , with  **$\rho$**  solution of

$$s = \text{mean}(\text{geom}(\rho)) + \frac{\lambda}{\mu} \rho$$

as  **$M, N$**  tends to  $+\infty$ ,

**$M/N \rightarrow s$**  the average number of bikes per station.

## A first result

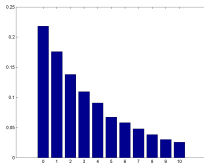
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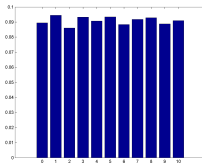
**$M/N \rightarrow s$**  the average number of bikes per station.

$$s < \frac{C}{2} + \frac{\lambda}{\mu}$$



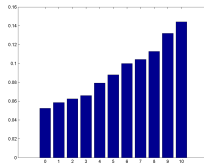
$\rho < 1$

$$s = \frac{C}{2} + \frac{\lambda}{\mu}$$



$\rho = 1$

$$s > \frac{C}{2} + \frac{\lambda}{\mu}$$



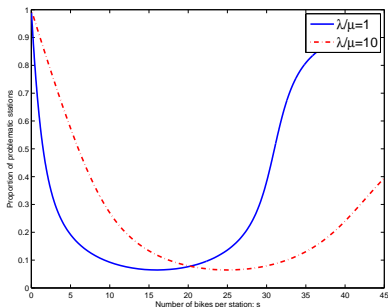
$\rho > 1$

# Plotting performance...

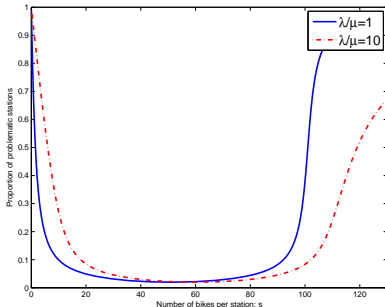
y-axis: limiting proportion of problematic stations

x-axis: number of bikes/station  $s$ .

$$\rho \mapsto (s(\rho), \text{geom}(\rho)(0) + \text{geom}(\rho)(C)).$$



(a)  $C = 30$ .



(b)  $C = 100$ .

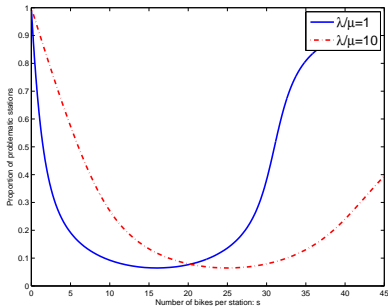
- ▶ minimal  $2/(C+1)$  for  $s = s_{opt} \stackrel{\text{def}}{=} \lambda/\mu + C/2$
- ▶ “flat” at  $s_{opt}$ .

## Plotting performance...

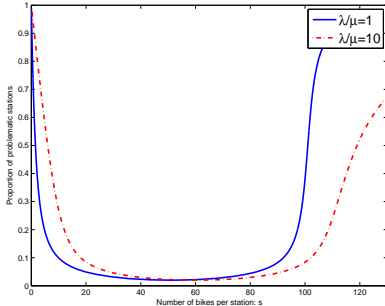
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(c)  $C = 30$ .



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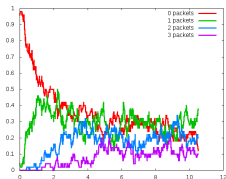
- ▶ minimal  $2/(C+1)$  for  $s = s_{opt} \stackrel{\text{def}}{=} \lambda/\mu + C/2$
- ▶ “flat” at  $s_{opt}$ .

Ex: for  $C = 30$ : at least 6.5% of problematic stations.

## Techniques: mean field limit

- Gives the steady state when  $N$  goes to infinity of

$$X_i = \frac{1}{N} \# \{ \text{stations with } i \text{ bikes} \}$$



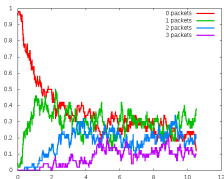
For fixed  $N$ ,  $X_i$  is a (reversible) Markov process

- Steady state (explicit)

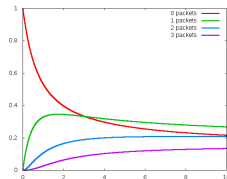
# Techniques: mean field limit

- Gives the steady state when  $N$  goes to infinity of

$$\mathbf{X}_i = \frac{1}{N} \# \{\text{stations with } i \text{ bikes}\} \rightarrow \frac{\rho^i}{Z}$$



$N \rightarrow \infty$   
→



For fixed  $N$ ,  $\mathbf{X}_i$  is a (reversible) Markov process

- Steady state (explicit)

System described by an ODE

- The ODE has a unique equilibrium point.
- Closed-form formula.

Key: proba. interpretation of the ODE

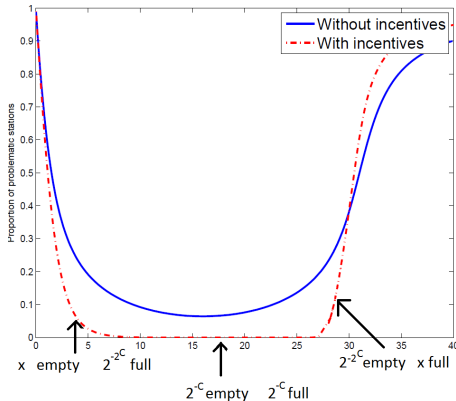
# Contributions on mean field

- ▶ mean field limit: explicit derivations
- ▶ CV of invariant measures via a Lyapunov function
- ▶ extension to some heterogeneity  
idea of proof: discrete/continuous



# Choice

- **Rule:** When returning, **choose at random two** stations and return to the **least loaded**.  
(Load balancing via power of choice)



- Improves performance from  **$1/C$**  to  **$2^{-C}$**
- True even if **incentives**=**just a proportion** choose

# Choice

- ▶ **Rule**: choose **two neighbors**
- ▶ **Simulations** (2D) show performance is  $\approx 2^{-C/2}$ .  
(recall: no geometry:  $2^{-C}$ , no incentive:  $1/C$ ).

# Conclusion

## Good understanding of the homogeneous case

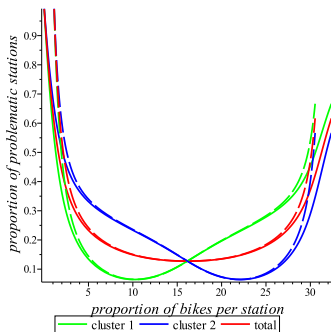
- ▶ **Poor performance:**  $1/C$  problematic stations (even in symmetric case)
- ▶ Choose among 2 neighbors helps a lot:  $2^{-c/2}$
- ▶ Avoid empty and full stations is **useless!**
- ▶ About **redistribution by trucks**
  - ▶ an simple efficient algorithm
  - ▶ a minimal value for truck rate of  $1/C$ : 0!

## Future work

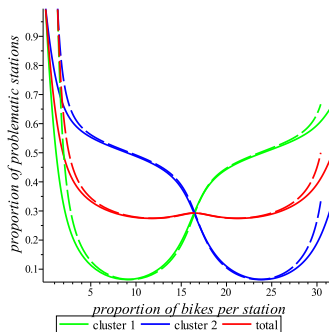
- ▶ mean field with geometry (local search)
- ▶ redistribution/incentives + heterogeneity: algorithms
- ▶ experimentation building a realistic model for Paris collaboration with Mairie de Paris

# Heterogeneity

From numerical simulations to theoretical results



(e)  $\lambda_1/\lambda_2 = 1.3$



(f)  $\lambda_1/\lambda_2 = 2$

Figure: Performance of the system

## A more complex system

- ▶ **heterogeneity**: housing vs working, uphill vs downhill,...
- ▶ **geometry** local search
- ▶ **time dependency**: daily period